

CSc 360

# Operating Systems

## Process Operations

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# Linux computers in ECS 242?!

- No longer remote accessible
  - due to the “flexilab” w/ hypervisors
  - both Linux and Win images
- Drop-in possible
  - see schedule
- Other alternatives
  - [linux.csc.uvic.ca](http://linux.csc.uvic.ca)
  - auto-load-balanced to multiple machines

LAB SCHEDULE  
ECS 242 - SUMMER 2015  
Lab is open for drop-in use during all unbooked times.

	Monday	Tuesday	Wednesday	Thursday	Friday
8am					
9am					
10am					ECS 242 801 10:00-11:00am
11am					ECS 242 802 11:00-12:00pm
12pm					
1pm			ECS 242 801 1:30-2:30pm		
2pm			ECS 242 801 2:30-3:30pm		
3pm			ECS 242 803 3:30-4:30pm		
4pm			ECS 242 804 4:30-5:30pm		
5pm					
6pm					
7pm					
8pm					

# Review: process

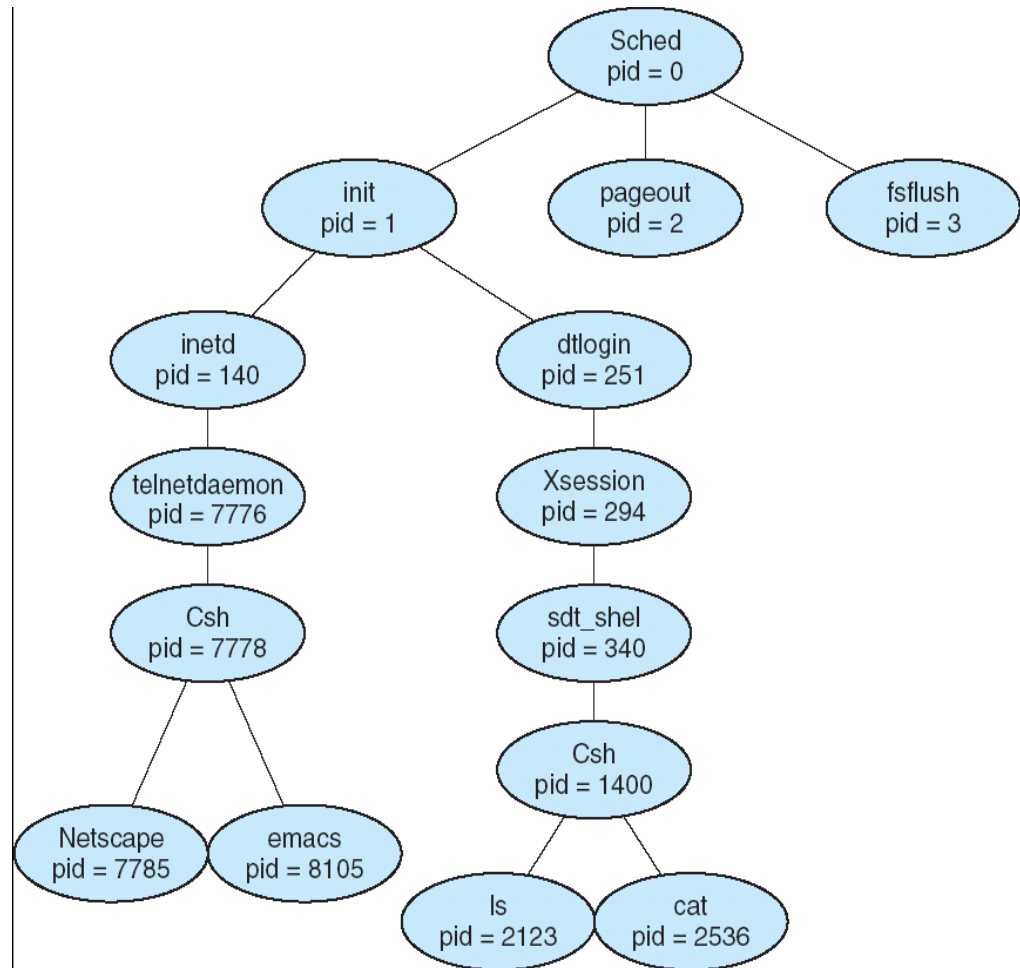
- Process vs program
  - active vs passive entity
- Process control block
  - context switching
- Process scheduling
  - long-term scheduling
  - short-term scheduling
  - medium-term scheduling

# Process creation

- Creating processes
  - parent process: create child processes
  - child process: created by its parent process
- Process tree
  - recursive parent-child relationship; why tree?
  - `/usr/bin/pstree`
- Process ID (PID) and Parent PID (PPID)
  - usually nonnegative integer

# Process tree

- sched (0)
  - init (1)
    - all user processes
  - pageout
    - memory
  - fsflush
    - file system



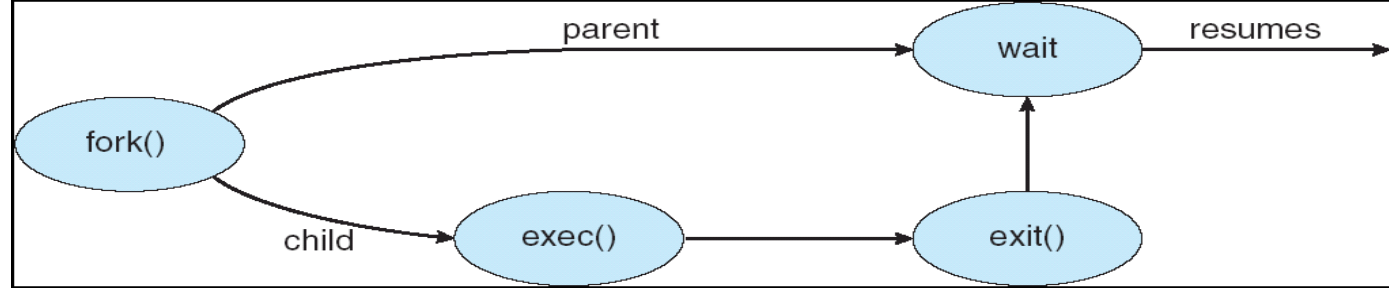
```
snmpd
sshd
  5*[sshd—sshd—sftp-server]
  sshd—sshd—bash—ssi
  sshd—sshd—bash—more
  sshd—sshd—bash—pstree
  sshd—sshd—csh—sftp-server
  sshd—sshd—bash
  sshd—sshd—bash—mutt
  sshd—sshd—tcsh—nano
  sshd—sshd—tcsh
syslogd
```

# Parent vs child processes

- Process: running program + resources
- Resource sharing: possible approaches
  - all shared
  - some shared (e.g., read-only code)
  - nothing shared\*
- Process execution: possible approaches
  - parent waits until child finishes
  - parent and child run concurrently\*

# fork(), exec\*(), wait()

- Create a child process: fork()
  - return code < 0: error (in “parent” process)
  - return code = 0: you’re in child process
  - return code > 0: you’re in parent process
    - return code = child’s PID
- Child process: load a new program
  - exec\*(): front-end for execve(file, arg, environ)
- Parent process: wait() and waitpid()



## Example

```
int main()
{
    Pid_t pid;
    /* fork another process */
    pid = fork();
    if (pid < 0) { /* error occurred */
        fprintf(stderr, "Fork Failed");
        exit(-1);
    }
    else if (pid == 0) { /* child process */
        execvp("/bin/ls", "ls", NULL);
    }
    else { /* parent process */
        /*parent will wait for the child to complete*/
        wait (NULL);
        printf ("Child Complete");
        exit(0);
    }
}
```



# Process termination

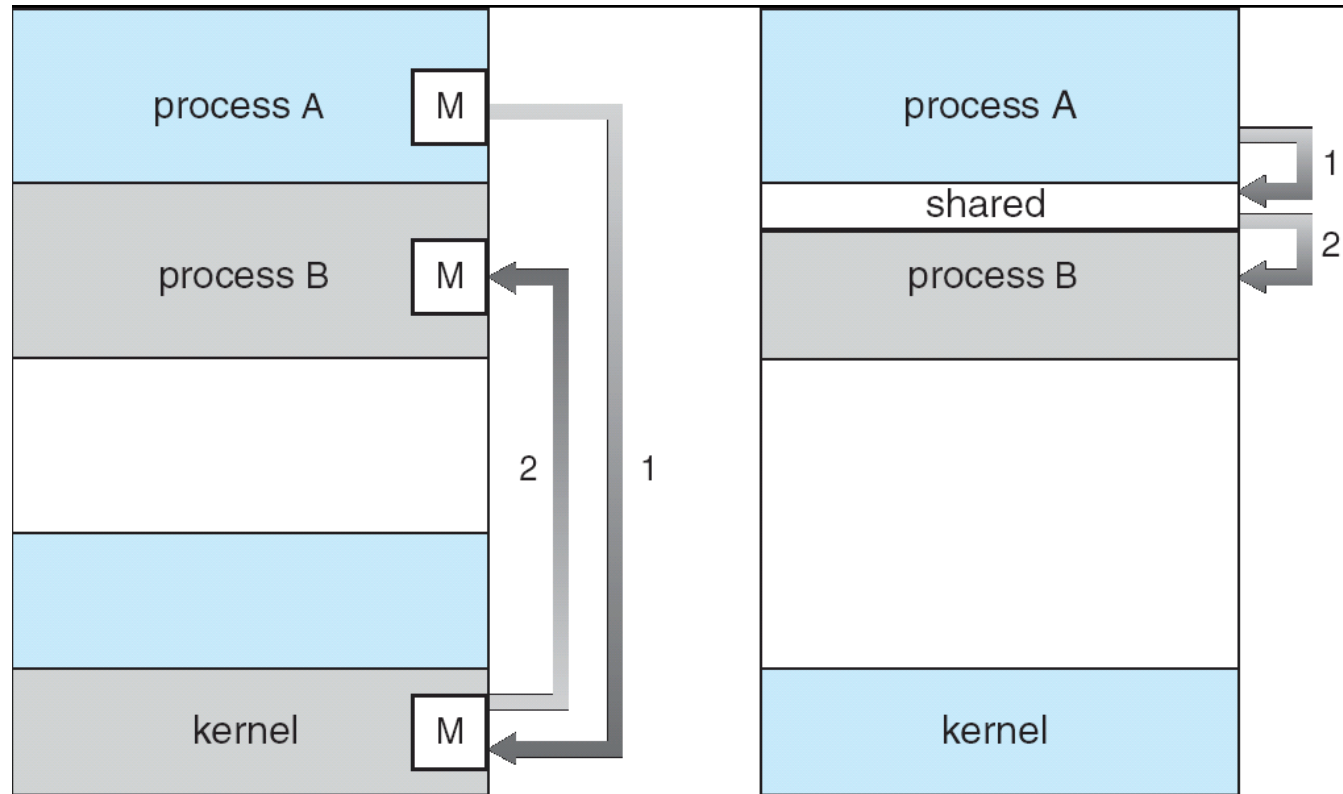
- Terminate itself: `exit()`
  - report status to parent process
  - release allocated resources
- Terminate child processes: `kill(pid, signal)`
  - actually send a signal to the child
    - child resource exceeded, child process no long needed, and so on
  - parent is exiting
    - cascading termination, or find another parent

# Process communication

- Independent process
  - standalone process
- Cooperating process
  - affected by or affecting other processes
    - sharing, parallel, modularity, convenience
- Process communication
  - shared memory
  - message passing

# Message passing vs shared memory

- Overhead
- Protection



(a)

(b)

# This lecture

- Process operations
  - process creation
    - process tree
  - process termination
  - the need for inter-process communication
- Explore further
  - /bin/ps, /usr/bin/top, /usr/bin/pstree
  - how does a child process find its parent's PID?

# Next lecture

- Inter-process communication
  - read OSC7 Chapter 3 (or OSC6 Chapter 4)